# Data Streams 

They're Everywhere<br>Magnetic Stripe as Lab Example

## Ubiquity of Data Streams

- We've seen $I^{2} C$
- and some intro to SPI and UART/RS-232
- Remote Controls (IR)
- pulses of infrared light
- Aircraft Transponders
- pulses of radio waves
- Cell Phone Data
- sophisticated modulation schemes, but still digital data
- Magnetic Stripe
- we'll use as a fun example in lab


## H-ITT Infrared Clickers

Old in-class clickers were IR: 0.5 ms pulse width; two similar packets back-to-back

Reverse-engineered transmission scheme: 1's fat; 0's skinny
end delimiter
Resulting bit-sequence for A signal (both packets) is:


ID is binary for 55573; that transmitter's permanent ID
A, B, D, C, E, * just counted up in binary: 1001 = A; 1010 = B, etc. checksum provides verification that data is correctly received

## What's with the Checksum?



Break data into chunks of 8 bits (bytes) and add up:

| 1001 |
| ---: |
| 00000000 |
| 11011001 |
| 00010101 |
| 11110111 |

Checksums provide a "sanity check" on the data integrity

## H-ITT RS-232 Datastream to Computer

E-button on H-ITT (first of two packets):


- Serial datastream looks a lot different
- this one allows many zeros or ones in a row
- delimiters (called start bit and stop bit) bracket 8-bit data (1 byte)
- in this case, 0's are positive voltage, 1's are negative (inverted; RS-232 std.)
- happens much faster than IR: in this case 19,200 bits/sec (baud)
- Packet breakdown:
- first packet: button number ( $5 \rightarrow$ E), with LSB first: 101000
- next three packets are ID, also LSB first within each
- last packet is checksum type of verification



## Stereo Remote Control

- Similar to H-ITT transmitters in principle:
- bursts of infrared light carrying digital information
- punctuated by delimiters so no long sequences of 1's or 0's
- Key differences:
- signal initiated by a WAKE UP! constant-on burst
- pattern that follows is repeated indefinitely until button is released
- I can never get fewer than three packets...
- packet is variable in length depending on button



## Sample patterns for data packet

| power |  |
| :---: | :---: |
| vol+ |  |
| vol- |  |
| 1 |  |
| 2 |  |
| 3 |  |
| 4 | ภ. |
| 5 | . -1 L- |
| 6 |  |
| 7 |  |

## A Different Code...

- The radio remote uses a different scheme:
- essentially nulls are $3 \times$ longer for 1 than for 0

ID part

data part

- in data part, least significant bit (LSB) is first
- here $0 \times 25$


## Aircraft Transponders at 1090 MHz



- Legacy of WWII Friend-or-Foe
- Bursts of RF power at 1090 MHz
- At left: 12-bit pattern
- 4 octal digits; 3 delimeters
- Below: newer data-rich

Even newer scheme at 978 MHz has more data, and error check scheme can correct several corrupted bytes in sequence
http://www.aircraft-avoid.com/ads-b-transition.html

- 56 or 112 bits
- can be lat/lon, etc.
- incl. parity check



## Magstripe Idea

- On magnetic stripe, N-S junctions eat their own magnetic flux lines, but N-N or S-S present external flux lines of opposite direction
- pattern of $\mathrm{N}-\mathrm{N}$ and $\mathrm{S}-\mathrm{S}$ creates + and transitions
- zero represented by long period
- one represented by short period
- zeros look fat; ones thin (sign irrelevant)
- two streams are produced from this:
- a data stream
- a clock
- data valid when clock high



## Magstripe Geometry



- Up to three tracks of data
- Tracks 1 and 3 typically higher-density (7-bit) alpha-numeric data
- Track 2 typically lower-density (5-bit) numeric data
- Track 2 used on almost every card; track 1 often, track 3 seldom


## Track 2 Character Code

| --Data |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| b1 | b2 | b3 | b4 | b5 | Character | Function |
| 0 | 0 | 0 | 0 | 1 | 0 ( 0 H ) | Data |
| 1 | 0 | 0 | 0 | 0 | 1 (1H) | " |
| 0 | 1 | 0 | 0 | 0 | 2 (2H) | " |
| 1 | 1 | 0 | 0 | 1 | 3 (3H) | " |
| 0 | 0 | 1 | 0 | 0 | 4 (4H) | " |
| 1 | 0 | 1 | 0 | 1 | 5 (5H) | " |
| 0 | 1 | 1 | 0 | 1 | 6 (6H) | " |
| 1 | 1 | 1 | 0 | 0 | 7 (7H) | " |
| 0 | 0 | 0 | 1 | 0 | 8 (8H) | " |
| 1 | 0 | 0 | 1 | 1 | 9 (9H) | " |
| 0 | 1 | 0 | 1 | 1 | : (AH) | Control |
| 1 | 1 | 0 | 1 | 0 | ; (BH) | Start Sentinel |
| 0 | 0 | 1 | 1 | 1 | $<$ ( CH ) | Control |
| 1 | 0 | 1 | 1 | 0 | $=(\mathrm{DH})$ | Field Separator |
| 0 | 1 | 1 | 1 | 0 | $>$ (EH) | Control |
| 1 | 1 | 1 | 1 | 1 | ? ( FH ) | End Sentinel |

## Track 2 Code Breakdown

- Five bits per character
- Last bit is Parity: ensures odd number of ones
- First four bits data: LSB first
- maps to: 0123456789 : ; <=>?
- numbers are direct binary mapping: $0110 \rightarrow 6$
- Control characters and formatting
- start sentinel is $11010 \rightarrow$;
- end sentinel is $11111 \rightarrow$ ?
- important that first bit of start is 1 so knows how to start slicing stream into chunks of 5


## Track 1/3

- Denser on stripe
- 7 bits per character
- odd parity (last bit, again)
- allows alpha-numeric set (6-bit data is 64 possibilities)
- ! "\#\$\%\&' ()*+,-./0123456789:; <=>?@ABCDEFGHIJKLMNOPQRSTUVWXYZ[\]^_
- "zeroth" character is a space, but can't see it here
- Start sentinel $1010001 \rightarrow$ \%
- note $5^{\text {th }}$ character (101 index; LSB first)
- End sentinel $1111100 \rightarrow$ ?
- note $31^{\text {st }}$ character (11111 index)


## Parity and LRC

- The parity catches single-bit errors
- but could get fooled by greater damage to data
- A longitudinal redundancy check (LRC) also employed
- one final character after end sentinel
- bit-wise running XOR combination of all prior chunks
- including start and end sentinels
- effectively 1 if odd number of 1 's in that bit position
- 0 doesn't alter running result; 1 flips from 0 to 1 or 1 to 0
- Extremely unlikely to get no parity errors AND match LRC


## Python/Pi Implementation

- After initial exploration on scope...
- manual study of bit patterns and example decode
- better to understand what computer needs to do
- Will let Pi take over
- Approach 1: polling
- constantly "ask" about digital values and have smarts to interpret
- Approach 2: interrupts
- wait for an edge (on the clock, or card-loaded) then sample data
- closer to what we do: look for clock pulse, check data there


## Interrupts on Pi GPIO

- Facilitated in standard RPi.GPIO library

```
import RPi.GPIO as GPIO
GPIO.setmode(GPIO.BCM)
GPIO.setup(some_pin, GPIO.IN, pull_up_down=GPIO.PUD_UP)
GPIO.wait_for_edge(some_pin, GPIO.FALLING, timeout=100)
```

- Using some_pin as stand-in for variable for BCM \#
- Pull-up input so high even if not asserted externally
- card reader signals are active low: idle high
- Wait for falling edge
- Set timeout (here 100 ms ) so it doesn't hang forever
- Pro-tip: first call to wait_for_edge takes 30-40 ms to release
- so call clock edge wait before getting into read loop
- otherwise miss initial data unless swipe speed is slow


## Program Flow

- Set up GPIO and constants (chunk size, character map)
- Wait for card swipe to start
- allow some number of seconds
- abort if no action
- Once card-loaded signal detected, begin collection
- on each clock edge, record data input channel value
- When done, break into chunks and process
- evaluate parity, character, and track LRC
- Report results in human-readable form


## Example Capture Code

```
loaded = False # starts not loaded (card not in)
seq = [] # list to hold sequence of ones and zeros
grace_ms = XXXX # decide how many milliseconds to allow
print "Swipe Card: you have %d seconds" % (grace_ms/1000.0)
beg = time.time() # grab a time in sec.
GPIO.wait_for_edge(XX, GPIO.FALLING, timeout=grace_ms) # card load edge
now = time.time() # grab post-load time
dt = now - beg # elapsed while waiting
if (dt > grace_ms/1000.0 - 0.1): # within 0.1 s of timeout
    print "Timed out. Cleaning up and exiting."
    GPIO.cleanup() # good form
    sys.exit() # exit program
else: # did not time out
    print "Card Load detected"
    loaded = True # register as legit
beg = time.time() # reset begin time
now = beg # start out now at beg
while ((now - beg) < XX and loaded): # give it some time
    GPIO.wait_for_edge(CLOCK_VAR_NAME,GPIO.FALLING,timeout=100) # caution delay
    bit = 1 - GPIO.input(DATA_VAR_NAME) # get data value; active low
    seq.append(bit) # append to running list
    if GPIO.input(CARD_LOADED_VAR_NAME): # replace name
        loaded = False # if high; no longer loaded
    now = time.time() # capture current time
```


## Process/Interpret Code

```
msg = ''
par = ''
pen = 0
lrc = 0
first = seq.index(1)
n_char = (len(seq) - first)/per
not_end = True
for ind in range(n_char):
    parcel = seq[first+per*ind:first+per*(ind+1)] # slice out one chunk
    n_ones = sum(parcel) # count up the ones in this chunk
    if n_ones % 2 == 0: # even number of ones (bad)
        par += 'X' # indicate bad
        if not_end: # still in valid sequence
            pen += 1
    else: # odd number of ones: parity good
        par += '.'
    strn = ''
    for val in parcel:
        strn += "%d" % val
    map_ind = int(strn[:per-1][::-1],2
    msg += charmap[map_ind]
    if not_end: # still have not seen end sentinel
        lrc ^= map_ind # accumulate LRC for valid data
    if (msg[-1] == '?' and ind > XX): # end sentinel for all tracks; after so many
        not_end = False # reached end of legit section
print "%s LRC = %s" % (msg,charmap[lrc])
print "%s; penalty = %d" % (par,pen)
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\section*{A Diagnostic Trick}
- If you need to sort out what's happening in your code, especially relative to the signal timing, insert a pulse to hardware:
```

GPIO.output(MONITOR_BCM,GPIO.HIGH)
time.sleep(0.0001)
GPIO.output(MONITOR_BCM,GPIO.LOW)

```
- Creates 0.1 ms pulse on some GPIO pin
- can then see where this comes, and if it happens at all

\section*{Reading}
- For magnetic stripe stuff, see:
- http://en.wikipedia.org/wiki/Magnetic stripe card
- http://money.howstuffworks.com/question503.htm
- http://stripesnoop.sourceforge.net/faq.html
- http://stripesnoop.sourceforge.net/devel/phrack37.txt```

